## **COMBINATORICA**

Bolyai Society – Springer-Verlag

# ON THE RELATION BETWEEN TWO MINOR-MONOTONE GRAPH PARAMETERS

#### R. PENDAVINGH

Received June 13, 1997

We prove that for each graph  $\mu(G) \leq \lambda(G) + 2$ , where  $\mu$  and  $\lambda$  are minor-monotone graph invariants introduced by Colin de Verdière [3] and van der Holst, Laurent, and Schrijver [5]. It is also shown that a graph G exists with  $\mu(G) < \lambda(G)$ . The graphs G with maximal planar complement and  $\mu(G) = |V(G)| - 4$ , characterised by Kotlov, Lovász, and Vempala, are shown to be forbidden minors for  $\{H \mid \mu(H) < |V(G)| - 4\}$ .

#### 1. Introduction

Given a graph G = (V, E) without loops or multiple edges, define  $\mathcal{O}_G$  as the collection of real-valued symmetric  $V \times V$  matrices  $M = (m_{ij})$  satisfying

- 1. if  $ij \in E$ , then  $m_{ij} < 0$ , and
- 2. if  $ij \notin E$  and  $i \neq j$ , then  $m_{ij} = 0$ .

There is no restriction on the diagonal entries. The elements of  $\mathcal{O}_G$  are sometimes called discrete Schrödinger operators.

A matrix  $M \in \mathcal{O}_G$  satisfies the *Strong Arnold Hypothesis*, SAH for short, if there is no nonzero symmetric matrix  $X = (x_{ij})$  such that MX = 0, and such that  $x_{ij} = 0$  whenever i = j or  $ij \in E$ .

By  $\lambda_i(M)$  we denote the *i*-th smallest eigenvalue of M. When G is connected and  $M \in \mathcal{O}_G$ , the Perron-Frobenius Theorem implies that for any eigenvector z of M:

z>0 or  $z<0 \iff z$  belongs to the smallest eigenvalue of M.

Hence, the multiplicity of  $\lambda_1(M)$  is 1. The parameter  $\mu(G)$  is defined as the largest corank of any matrix  $M \in \mathcal{O}_G$  with exactly one negative eigenvalue, satisfying the SAH.

This definition is due to Colin de Verdière [3]. Equivalently,  $\mu(G)$  is the maximum multiplicity of  $\lambda_2(M)$ , where M ranges over all  $M \in \mathcal{O}_G$  satisfying (a more general formulation of) the SAH. See [7] for a survey of results concerning  $\mu$ .

Mathematics Subject Classification (1991): 05C

If G = (V, E) is a graph and  $S \subseteq V$ , let G[S] denote the subgraph of G induced by S. Given  $x \in \mathbb{R}^V$ , the support of x is  $\operatorname{supp}(x) := \{v \in V \mid x_v \neq 0\}$ . Furthermore,  $\operatorname{supp}_+(x) := \{v \in V \mid x_v > 0\}$  is the postive support, and  $\operatorname{supp}_-(x) := \{v \in V \mid x_v < 0\}$  is the negative support of a vector x.

A linear subspace  $L \subseteq \mathbb{R}^V$  is said to be a valid representation of G when for each nonzero  $x \in L$ , one has

- 1.  $\operatorname{supp}_{+}(x) \neq \emptyset$ , and
- 2.  $G[\operatorname{supp}_{+}(x)]$  is connected.

Van der Holst, Laurent, and Schrijver [5] defined  $\lambda(G)$  as

$$\lambda(G) := \max\{\dim(L) \mid L \text{ is a valid representation of } G\}.$$

When G=(V,E) is a graph and  $S\subseteq V$ , we denote the set of neighbors of S in G by  $N_G(S)$ , i.e.  $N_G(S):=\{v\in V\setminus S\mid \exists w\in S, vw\in E\}$ . When  $M\in \mathcal{O}_G,\ M_S$  denotes the restriction of M to the rows and columns indexed by S. Given a vector  $x\in \mathbb{R}^V,\ x_S\in \mathbb{R}^S$  denotes the restriction of x to the positions with indices in S. By extension with zeros a vector  $x\in R^S$  may still be regarded as an element of  $R^V$ , and as such vectors restricted to different subsets of V may be added.

A graph H is a *subgraph* of a graph G, denoted  $H \subseteq G$ , if H can be obtained by removing vertices and deleting edges from G. When H can be obtained from a subgraph of G by contracting edges, H is a *minor* of G, which we denote by  $H \subseteq G$ .

Both  $\lambda$  and  $\mu$  are minor-monotone, i.e.

if H is a minor of G then  $\lambda(H) \leq \lambda(G)$  and  $\mu(H) \leq \mu(G)$ .

The short proof of the minor-monotony of  $\lambda$  appears in [5]. The SAH plays an essential role in proving that  $\mu$  is minor-monotone. Colin de Verdière's original proof can be found in [3].

A graph G = (V, E) is a *clique sum* of  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  if  $V = V_1 \cup V_2, E = E_1 \cup E_2$  and  $V_1 \cap V_2$  induces a clique in both  $G_1$  and  $G_2$ . If G is a clique sum of  $G_1$  and  $G_2$ , then  $\lambda(G) \geq \lambda(G_i)$  and  $\mu(G) \geq \mu(G_i)$  as  $G_1$  and  $G_2$  are both subgraphs of G. By work of Van der Holst, Laurent, and Schrijver [5], we know that

$$\lambda(G) = \max\{\lambda(G_1), \lambda(G_2)\}.$$

Van der Holst, Lovász, and Schrijver [7] showed that either

- 1.  $\mu(G) \le \max\{\mu(G_1), \mu(G_2)\}\$ , or
- 2.  $\mu(G) = \max\{\mu(G_1), \mu(G_2)\} + 1$ , and  $\mu(G_1) = \mu(G_2)$ .

It follows inductively that if G is a clique sum of more than two graphs  $G_1, \ldots, G_k$ , then  $\mu(G) \leq \max\{\mu(G_1), \ldots, \mu(G_k)\} + 1$ .

A  $Y\Delta$ -operation on G is removing a vertex v of degree 3 from G, and then adding a triangle on the former neighbors of v. The converse is called a  $\Delta Y$ -operation. Bacher and Colin de Verdière show [1]: if H is obtained from G by a  $Y\Delta$ -operation, then  $\mu(H) \leq \mu(G)$ . Also, if G is obtained from H by a  $\Delta Y$ -operation, G is a subgraph of a clique sum of H and  $K_4$ . Hence, if  $\mu(H) \geq 4 > \mu(K_4)$ , we have  $\mu(G) = \mu(H)$ .

Given any set of graphs  $\mathcal C$  closed under taking minors, define the forbidden minors for  $\mathcal C$  as

$$F(\mathcal{C}) := \{ G \mid G \notin \mathcal{C}, \forall H < G \mid H \in \mathcal{C} \}.$$

Such a set of forbidden minors is finite, by the Robertson-Seymour graph minor Theorem. Clearly, membership of  $\mathcal{C}$  can be characterized by

$$H \in \mathcal{C} \Leftrightarrow H$$
 has no element of  $F(\mathcal{C})$  as a minor.

Since  $\mu$  is minor-monotone, the set  $\{G | \mu(G) < k\}$  is closed under taking minors, for any k. The following forbidden minor characterizations of  $\mu < k$  are known:

- 1.  $\mu(G) < 1 \iff G$  has no  $\overline{K_2}$ -minor,
- 2.  $\mu(G) < 2 \iff G$  has no  $K_3$  or  $K_{1,3}$ -minor,
- 3.  $\mu(G) < 3 \iff G$  has no  $K_4$  or  $K_{2,3}$ -minor,
- 4.  $\mu(G) < 4 \iff G$  has no  $K_5$  or  $K_{3,3}$ -minor,
- 5.  $\mu(G) < 5 \iff G$  has no minor in the Petersen family.

In each of these statements, the ' $\Longrightarrow$ ' part is relatively easy to verify: it suffices to compute  $\mu$  for the graphs mentioned on the right. All but the last of these results are due to Colin de Verdière [3]. For the characterisation of  $\mu < 4$  he used Kuratowski's Theorem, that a graph without  $K_5$ - or  $K_{3,3}$ -minor is planar.

The characterisation of  $\mu < 5$  is due to Lovász and Schrijver [9]. They show that so-called 'flat' graphs have  $\mu < 5$ . By a Theorem of Robertson, Seymour, and Thomas [10], a graph without a minor in the Petersen family is a flat graph. The Petersen family is the set of 7 graphs that can be obtained from  $K_6$  by any series of  $\Delta Y$ - and  $Y\Delta$ -operations. The Petersen graph is in the Petersen family.

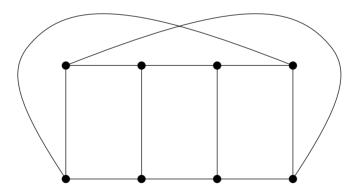


Figure 1. The graph  $V_8$ 

The corresponding results on  $\lambda$  are [5]:

- 1.  $\lambda(G) < 1 \iff G$  has no  $\overline{K_2}$ -minor,
- 2.  $\lambda(G) < 2 \iff G \text{ has no } K_3\text{-minor},$
- 3.  $\lambda(G) < 3 \iff G$  has no  $K_4$ -minor,

4.  $\lambda(G) < 4 \iff G$  has no  $K_5$ - or  $V_8$ -minor.

The graph  $V_8$  is shown in Figure 1.

It was speculated in [9] that the following might be true:

a graph G satisfies  $\lambda(G) \le t$  if and only if G is a subgraph of a clique sum of graphs H with  $\mu(H) \le t$ .

The 'only if' part implies that for any graph G, we have  $\mu(G) \leq \lambda(G) + 1$ . The 'if' part implies that  $\mu(G) \geq \lambda(G)$  for all graphs G.

In the next section it is shown that for any graph G, we have  $\mu(G) \leq \lambda(G) + 2$ . In section 4, we find a graph G with  $\mu(G) < \lambda(G)$ .

#### 2. A relation between $\mu$ and $\lambda$

The starting point of the present investigation is the following lemma. Its proof appears in [7].

**Lemma 1.** (van der Holst) Let G be a connected graph and let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue. Let  $x \in \ker M$  be such that  $G[\operatorname{supp}_+(x)]$  is disconnected. Then there is no edge between  $\operatorname{supp}_+(x)$  and  $\operatorname{supp}_-(x)$ , and for each component C of  $G[\operatorname{supp}(x)]$ ,  $N_G(C) = N_G(\operatorname{supp}(x))$ .

Given a graph G = (V, E), call a vector  $x \in \mathbb{R}^V$  broken when  $G[\operatorname{supp}(x)]$  has at least 3 components. Observe that for any connected graph G, if  $M \in \mathcal{O}_G$  has exactly one negative eigenvalue, and x is a nonzero vector in  $\ker(M)$ , then

- 1.  $\operatorname{supp}_+(x) \neq \emptyset$ , and
- 2. if x is not broken, then  $G[\operatorname{supp}_{+}(x)]$  is connected.

(1. holds as x is orthogonal to an eigenvector z>0 of M corresponding to  $\lambda_1(M)$ , and 2. is a consequence of Lemma 1.)

So each subspace of  $\ker(M)$  avoiding broken vectors is a valid representation of G.

When  $M \in \mathcal{O}_G$  has exactly one negative eigenvalue with corresponding eigenvector z, then Rayleigh's Theorem implies

$$x \perp z$$
 and  $x^T M x = 0 \iff x \in \ker(M)$ .

Using this fact van der Holst, Lovász, and Schrijver [6, 7] show:

**Lemma 2.** Let G be a connected graph and  $M \in \mathcal{O}_G$  with exactly one negative eigenvalue. Let  $S \subseteq V(G)$ , and  $C_1, \ldots, C_k$ , the components of G - S. Then  $\lambda_1(M_{C_i}) < 0$  implies  $\lambda_1(M_{C_i}) > 0$  for all  $j \neq i$ .

With the techniques that were used to prove Lemma 1 and Lemma 2, it is possible to show the following:

**Lemma 3.** Let G be a connected graph and let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue. Let  $x \in \ker(M)$ , and set

$$D := \{ y \in \ker(M) \mid \operatorname{supp}(y) \subseteq \operatorname{supp}(x) \}.$$

If  $G[\operatorname{supp}(x)]$  is disconnected, it has exactly  $\dim(D) + 1$  connected components. If in addition M satisfies the SAH, then  $\dim(D) \leq 2$ .

**Proof.** Let  $G[\operatorname{supp}(x)]$  be disconnected, and let  $C_1, \ldots, C_k$  be its connected components. Mx=0 implies  $M_{C_i}x_{C_i}=0$ , and hence  $\lambda_1(M_{C_i})\leq 0$  for each  $i=1,\ldots,k$ .

By Lemma 2,  $\lambda_1(M_{C_i}) < 0$  would imply  $\lambda_1(M_{C_j}) > 0$  for all  $j \neq i$ , and thus we have  $\lambda_1(M_{C_i}) = 0$  for each i = 1, ..., k. By the Perron-Frobenius Theorem applied to each  $M_{C_i}$  we have either  $x_{C_i} < 0$  or  $x_{C_i} > 0$ . Hence if  $y \in D$  then there exists  $\alpha_i \in \mathbb{R}$  such that  $y_{C_i} = \alpha_i x_{C_i}$ , since  $M_{C_i} y_{C_i} = 0$ . Let z be an eigenvector corresponding to the smallest eigenvalue of M. If  $y \in \ker(M)$  then  $z^T y = 0$ , hence D is contained in

$$D' := \{ y \in \mathbb{R}^V \mid y = \sum \alpha_i x_{C_i}, \alpha_i \in \mathbb{R}, z^T y = 0 \}.$$

On the other hand, suppose  $y \in D'$ . Then, since  $y^T M y = \sum \alpha_i^2 x_{C_i}^T M_{C_i} x_{C_i} = 0$  and  $z^T y = 0$  we have  $y \in \ker(M)$  by Rayleigh's Theorem, and certainly  $\sup(y) \subseteq \sup(x)$ , so D = D'. Since  $z^T y = 0 \Leftrightarrow \sum \alpha_i (z_{C_i}^T x_{C_i}) = 0$ , and  $z_{C_i}^T x_{C_i} \neq 0$  for all i, it follows that  $\dim(D) = \dim(D') = k - 1$ .

Now assume M satisfies the SAH, and suppose for a contradiction that  $G[\operatorname{supp}(x)]$  has more than 3 components. Clearly there exist  $s,t\in D$  such that  $\operatorname{supp}(s)=C_1\cup C_2$  and  $\operatorname{supp}(t)=C_3\cup C_4$ . But then  $X=ts^T+st^T$  is a symmetric matrix such that MX=0 and

$$x_{ij} \neq 0 \Rightarrow (i \in C_1 \cup C_2, j \in C_3 \cup C_4 \text{ or vice versa}) \Rightarrow ij \notin E(G),$$
 contradicting the SAH.

**Theorem 1.** For all connected graphs G,  $\mu(G) \leq \lambda(G) + 2$ .

**Proof.** Let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue and satisfy the SAH, with  $\operatorname{corank}(M) = \mu(G)$ . By Lemma 3, all broken vectors in  $\ker(M)$  are contained in finitely many 2-dimensional subspaces of  $\ker(M)$ . Then there exists a subspace  $L \subseteq \ker(M)$  of dimension  $\operatorname{corank}(M) - 2$  that has no nonzero vector in common with any of these subspaces. But then L is a valid representation of G, and  $\lambda(G) \ge \dim(L) = \mu(G) - 2$ .

The question remains whether this bound is sharp. We do not know any graph G having  $\mu(G)=\lambda(G)+2$ , so it may be true that  $\mu(G)\leq \lambda(G)+1$  for all G. Indeed, the latter bound holds for graphs with  $\mu(G)\leq 5$ , as one verifies knowing the forbidden minor characterizations of  $\{G\mid \mu(G)< k\}$  for  $k=1,\ldots 5$ . Also,  $\mu(K_{k,l})\leq \min\{k,l\}+1=\lambda(K_{k,l})+1$  [7]. It is certainly not true that  $\mu(G)\leq \lambda(G)$  for all G. For example,  $\mu(K_{1,3})=2=\lambda(K_{1,3})+1$  and  $\mu(P)=5=\lambda(P)+1$  for the Petersen graph P.

#### 3. 3-regular graphs and graphs on a surface

For 3-regular and claw-free graphs, we obtain somewhat better bounds than the  $\mu \leq \lambda + 2$  of Theorem 1.

**Lemma 4.** Let G be a connected graph and let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue. Suppose there is an edge  $e = uv \in E(G)$  such that

$$e \cap \operatorname{supp}(x) \neq \emptyset$$
 for all broken  $x \in \ker(M)$ .

Then  $\operatorname{corank}(M) \leq \lambda(G-e) + 1$ .

**Proof.** We show that the space  $L := \{y \in \ker(M) \mid y_u + y_v = 0\}$  is a valid representation of G-e. On the contrary, suppose that some  $y \in L$  is broken. Since for each component  $C_i$  of  $G[\operatorname{supp}(y)]$  either  $y_{C_i} > 0$  or  $y_{C_i} < 0$ , and by assumption  $e \cap \operatorname{supp}(y) \neq \emptyset$ , e has an endpoint in two different components of  $G[\operatorname{supp}(y)]$ , a contradiction with Lemma 1. So y is not broken and hence  $G[\operatorname{supp}_+(y)]$  is connected. Because  $y_u + y_v = 0$ , e does not have both endpoints in  $\operatorname{supp}_+(y)$ . So even  $(G-e)[\operatorname{supp}_+(y)]$  is connected. Hence, L is a valid representation of G-e, and  $\lambda(G-e) \geq \dim(L) \geq \operatorname{corank}(M) - 1$ .

**Lemma 5.** Let G be a connected, 3-regular graph and let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue. If  $x \in \ker(M)$  is a broken vector,  $G - \operatorname{supp}(x)$  has no edges.

**Proof.** By Lemma 1 we have  $N_G(C) = N_G(\operatorname{supp}(x))$  for each component C of  $G[\operatorname{supp}(x)]$ . Since G is 3-regular, a vertex  $v \not\in \operatorname{supp}(x)$  has either  $N_G(v) \subseteq \operatorname{supp}(x)$  or  $N_G(v) \cap \operatorname{supp}(x) = \emptyset$ . By the connectedness of G, if  $V \setminus \operatorname{supp}(x) \setminus N_G(\operatorname{supp}(x))$  is nonempty it is connected to  $N_G(\operatorname{supp}(x))$ , but then there exists a vertex  $v \in N_G(\operatorname{supp}(x))$  that is connected to  $V \setminus \operatorname{supp}(x) \setminus N_G(\operatorname{supp}(x))$ , a contradiction. Hence  $V \setminus \operatorname{supp}(x) = N_G(\operatorname{supp}(x))$  and  $N_G(v) \subseteq \operatorname{supp}(x)$  for all  $v \not\in \operatorname{supp}(x)$ .

**Theorem 2.** For any connected 3-regular graph G,  $\mu(G) \leq \lambda(G) + 1$ .

**Proof.** Let  $M \in \mathcal{O}_G$  have exactly one eigenvalue and  $\operatorname{corank}(M) = \mu(G)$ . Let  $e \in E(G)$  be any edge. By Lemma 5,  $e \cap \operatorname{supp}(x) \neq \emptyset$  for all broken  $x \in \ker(M)$ . By Lemma 4,  $\mu(G) = \operatorname{corank}(M) \leq \lambda(G - e) + 1 \leq \lambda(G) + 1$ .

For example, the Petersen graph P has  $\mu(P) = 5$  and  $\lambda(P) = 4$ , and when an arbitrary edge is removed from P the result is a subdivision of  $V_8$  (see figure 1). The above proof yields a construction of a 4-dimensional valid representation of  $V_8$ , given a matrix  $M \in \mathcal{O}_P$  with exactly one eigenvalue and  $\operatorname{corank}(M) = 5$ . It is not necessary that the SAH holds for such a matrix M.

**Theorem 3.** If G is a connected claw-free graph, then  $\mu(G) \leq \lambda(G)$ .

**Proof.** Suppose  $M \in \mathcal{O}_G$  with exactly one eigenvalue and  $\operatorname{corank}(M) = \mu(G)$ . Suppose  $x \in \ker(M)$  is broken. As G is connected, there exists  $v \in N_G(\operatorname{supp}(x))$ . Then, by Lemma 1, v has neighbors in each of at least 3 components of  $\operatorname{supp}(x)$ .

This is a contradiction with the assumption that G is claw-free. Hence  $\ker(M)$  is a valid representation of G of dimension  $\mu(G)$ .

**Theorem 4.** Given any surface S,

$$\max\{\mu(H) \mid H \text{ embeds in } S\} \leq \max\{\lambda(H) \mid H \text{ embeds in } S\}.$$

**Proof.** Let G attain the maximum in  $\max\{\mu(H) \mid H \text{ embeds in } S\}$ . We will construct a clawfree graph G' that has G as a minor and is embedded in S. By minor-monotony and Theorem 3 we then have  $\mu(G) \leq \mu(G') \leq \lambda(G')$ , and the Theorem will follow.

We may assume that all vertices of G have degree at least 3. To obtain G' from G, first split vertices such that embeddability in S is preserved, until each vertex has degree 3. Next, replace each edge by a path of length 2. Finally, add edges connecting each pair of neighbors of a vertex of degree 3. The resulting graph G' has G as a minor by construction, is claw-free and is embedded in S.

### 4. An upper bound on $\mu$ , and a counterexample

**Theorem 5.** If G = (V, E) is any connected graph, then either

$$|E| \ge \mu(G)(\mu(G) + 1)/2$$
 or  $G = K_{3,3}$ .

**Proof.** Suppose there exists some  $G \neq K_{3,3}$  with  $|E| < \mu(G)(\mu(G) + 1)/2$ . Then  $\mu(G) > 4$ , and we may assume that G is triangle-free, as applying  $\Delta Y$  operations does not violate the condition that  $|E| < \mu(G)(\mu(G) + 1)/2$ .

Let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue and satisfy the SAH, with  $\operatorname{corank}(M) = \mu(G)$ .

Suppose that some diagonal entry of M is nonzero, say  $m_{11} \neq 0$ . Let  $F \subseteq E$  be the edges of a spanning tree of G. The linear space of matrices

$$\mathcal{X} := \{X \mid X \text{ symmetric and } MX = 0\}$$

has dimension  $\dim(\mathcal{X}) = \mu(G)(\mu(G)+1)/2$ . Consider the subspace

$$\mathcal{X}' := \{ X \in \mathcal{X} \mid \forall ij \in E \setminus F \ x_{ij} = 0, \forall i \neq 1 \ x_{ii} = 0 \}.$$

Let  $X \in \mathcal{X}'$ , and set  $F' := \{ij \in E \mid x_{ij} \neq 0\}$ . As  $F' \subseteq F$ , the graph G' = (V, F') is a forest. If  $F' \neq \emptyset$ , then there exists some vertex  $i \neq 1$  of degree 1 in G'. This contradicts the fact that  $x_{ii} = 0$  and  $m_i^T x_i = 0$ . Hence,  $x_{ij} = 0$  for all  $ij \in E$ . Similarly,  $m_1^T x_1 = 0$  and  $m_{11} \neq 0$  imply  $x_{11} = 0$ . So also  $x_{ii} = 0$  for all  $i \in V$ .

Since the SAH holds for M we have X=0. It follows that  $\mathcal{X}'$  contains no nonzero elements. From this, and the definition of  $\mathcal{X}'$  it follows that

$$0 = \dim(\mathcal{X}') \ge \dim(\mathcal{X}) - |E \setminus F| - (|V| - 1)$$
$$= \mu(G)(\mu(G) + 1)/2 - |E| > 0,$$

a contradiction. So  $m_{ii} = 0$  for all  $i \in V$ .

Let z>0 be an eigenvector belonging to the smallest eigenvalue of M. If  $ij \notin E$ , then the vector u defined by  $u_i=z_j$ ,  $u_j=-z_i$  and  $u_k=0$  when  $k\neq i,j$  has  $z^Tu=0$  and  $u^TMu=0$ . Hence,  $u\in \ker(M)$  by Rayleigh's Theorem. By Lemma 1 applied to u, it follows that N(i)=N(j). Hence, G is complete multipartite. If C is a coclique of G, then we similarly find an  $x\in \ker(M)$  with  $\operatorname{supp}(x)=C$ . Hence  $|C|\leq 3$  by Lemma 3. Since G is also triangle-free, G is an induced subgraph of  $K_{3,3}$ , contradicting the fact that  $\mu(G)>4$ .

For any k, the complete graph  $K_{k+1}$  is a forbidden minor for  $\{H \mid \mu(H) < k\}$ . Also, if k > 4 then each graph obtained from  $K_{k+1}$  by any series of  $\Delta Y$ - and  $Y\Delta$ -operations is a forbidden minor, and all these graphs have the same number of edges as  $K_{k+1}$ . In general, not every forbidden minor is obtained this way, but we do have:

**Corollary 5.1.** If G is a forbidden minor for  $\{H | \mu(H) < k\}$ , then G has at least as many edges as the complete graph  $K_{k+1}$ , or  $G = K_{3,3}$ .

Kotlov, Lovász, and Vempala [8] characterise the graphs G whose complement  $\overline{G}$  is a maximal planar graph and for which  $\mu(G) \ge |V(G)| - 4$ .

**Corollary 5.2.** If  $\overline{G}$  is a maximal planar graph and  $\mu(G) = |V(G)| - 4$ , then G is a forbidden minor for  $\{H \mid \mu(H) < |V(G)| - 4\}$ .

**Proof.** If  $\overline{G}$  is a maximal planar graph, then  $|E(\overline{G})| = 3|V(G)| - 6$ . Hence, |E(G)| = (|V(G)| - 4)(|V(G)| - 3)/2. Any proper minor H of G has strictly fewer edges than G, and hence by Theorem 5 we have  $\mu(H) < |V(G)| - 4$ .

For example, we have  $\mu(\overline{I}) = |V(\overline{I})| - 4 = 8$  where I denotes the Icosahedron. So  $\overline{I}$  is a forbidden minor for  $\{H | \mu(H) < 8\}$ , as is  $K_9$ . One cannot obtain  $\overline{I}$  from  $K_9$  by a series of  $\Delta Y$ - and  $Y\Delta$ -operations (shown by computer).

A generalised dodecagon of order (1,2) or GD(1,2) is a graph G with the following properties:

- 1. G is 3-regular,
- 2. G has diameter 6,
- 3. for each vertex  $v \in V$  and each i = 1, ..., 5: if  $u \in V$  is at distance i from v, then u has 2 neighbors at distance i+1 from v and 1 neighbor at distance i-1 from v, and

4. for each vertex  $v \in V$ : if  $u \in V$  has distance 6 from v, then it has 3 neighbors at distance 5 from v.

In other words, a GD(1,2) is a distance-regular graph with intersection array  $\{3,2,2,2,2,2;1,1,1,1,1,3\}$ . A GD(1,2) has 126 vertices, and 189 edges. There exists a unique generalised dodecagon of order (1,2) [2], known as Tutte's 12-Cage. We will denote this graph by T.

**Theorem 6.**  $\mu(T) \le 18 < 20 \le \lambda(T)$ .

**Proof.** By Theorem 5,  $\mu(T) \leq 18$  as |E(T)| = 189 and  $T \neq K_{3,3}$ . The second-largest eigenvalue  $\theta$  of the adjacency matrix A of T has multiplicity 21 [2, p. 416]. Hence  $M := \theta I - A \in \mathcal{O}_T$  has exactly one negative eigenvalue, and  $\operatorname{corank}(M) = 21$ . As T is 3-regular,  $\lambda(T) \geq 20$  by the proof of Theorem 2.

#### 5. Lower bounds for $\lambda$

The following theorem is due to Kotlov, Lovász, and Vempala [8]:

**Theorem 7.** For every graph G,

- 1. if  $\overline{G}$  is a disjoint union of paths, then  $\mu(G) \ge |V(G)| 3$ ,
- 2. if  $\overline{G}$  is outerplanar, then  $\mu(G) \ge |V(G)| 4$ ,
- 3. if  $\overline{G}$  is planar, then  $\mu(G) \ge |V(G)| 5$ .

These bounds on  $\mu$  can be extended to  $\lambda$  as follows:

**Theorem 8.** For every graph G,

- 1. if  $\overline{G}$  is a disjoint union of paths, then  $\lambda(G) \ge |V(G)| 3$ ,
- 2. if  $\overline{G}$  is outerplanar, then  $\lambda(G) \ge |V(G)| 4$ ,
- 3. if  $\overline{G}$  is planar and  $K_{2,2,2} \not\subseteq \overline{G}$ , then  $\lambda(G) \ge |V(G)| 5$ .
- **Proof.** 1. When  $\overline{G}$  is a disjoint union of paths,  $\mu(G) \geq |V(G)| 3$ . Also, G is claw-free as  $\overline{G}$  has no triangles. By Theorem 3,  $\mu(G) \leq \lambda(G)$ .
- 2. Observe that if  $\overline{H}$  is a subgraph of  $\overline{G}$ , then  $|V(H)| \lambda(H) \leq |V(G)| \lambda(G)$ . We may therefore assume that  $\overline{G}$  is maximally outerplanar. So  $\overline{G}$  is 2-connected, contains no  $K_{2,3}$  or  $K_4$ -minor, and has  $\mu(G) \geq |V(G)| 4$ .
- Let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue and  $\operatorname{corank}(M) = \mu(G)$ . Suppose  $x \in \ker(M)$  is broken. Let  $C_1, C_2, C_3$  be components of  $G[\operatorname{supp}(x)]$ , indexed such that  $|C_1| \ge |C_2| \ge |C_3|$ .

Since  $K_{|C_1|,|C_2|,|C_3|}$  is a subgraph of  $\overline{G}[\text{supp}(x)]$ , and  $K_{2,3}$  is not a subgraph of  $\overline{G}$ , either  $(|C_1|,|C_2|,|C_3|) = (1,1,1)$  or  $(|C_1|,|C_2|,|C_3|) = (2,1,1)$ . Similarly,

 $V \setminus \operatorname{supp}(x) = N_G(\operatorname{supp}(x))$  as  $K_4$  is not a subgraph of  $\overline{G}$ . Hence, if  $C_i$  is a singleton this implies that  $C_i$  is not connected to  $V \setminus \operatorname{supp}(x)$  in  $\overline{G}$ . As a maximally outerplanar graph is 2-connected,  $(|C_1|, |C_2|, |C_3|) = (1, 1, 1)$  entails  $\operatorname{supp}(x) = V$ , and we are done. When  $(|C_1|, |C_2|, |C_3|) = (2, 1, 1)$ ,  $v \notin \operatorname{supp}(x)$  must be connected to  $C_1$  by 2 vertex-disjoint paths in  $\overline{G}$ . This implies  $\overline{G} \geq K_{2,3}$ , a contradiction. So  $\operatorname{supp}(x) = V$  and we are done.

Since  $\ker(M)$  contains no broken vectors, it is a valid representation of G.

3. We may assume that  $\overline{G}$  is maximally planar, no vertex has degree 4 in  $\overline{G}$ , and  $K_{2,2,2} \not\subseteq \overline{G}$ , by the following argument. If a graph is not maximally planar and does not contain a  $K_{2,2,2}$ -subgraph, it is always possible to add an edge, keeping planar and not introducing a  $K_{2,2,2}$ -subgraph. Furthermore, when each triangular face of a maximally planar graph is subdivided by a vertex, the resulting graph contains no vertex of degree 4. Since  $\overline{G}$  is maximally planar,  $\overline{G}$  is 3-connected and contains no  $K_{3,3}$ - or  $K_5$ -minor.

Let  $M \in \mathcal{O}_G$  have exactly one negative eigenvalue and corank $(M) = \mu(G)$ , and let M satisfy the SAH. Suppose  $x \in \ker(M)$  is broken. Let  $C_1, C_2, C_3$  be the components of G[supp(x)], indexed such that  $|C_1| \ge |C_2| \ge |C_3|$ .

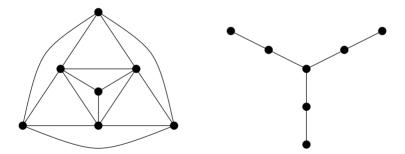


Figure 2.  $\overline{G}$  (left) is max. planar, and G (right) is a tree.

Since  $K_{3,3} \not\subseteq \overline{G}[\operatorname{supp}(x)]$  we have that if  $|C_1| \ge 3$  then  $|C_1| = |C_2| = 1$ . Hence there are the following cases:

Case 1.  $(|C_1|, |C_2|, |C_3|) = (1, 1, 1)$ :  $v \in N_G(\text{supp}(x))$  and 3-connectivity imply  $|V \setminus \text{supp}(x) \setminus N_G(\text{supp}(x))| \ge 3$ , as v must be 3-connected in  $\overline{G}$  to the vertex in  $C_3$ . Hence  $K_{3,3} \subseteq \overline{G} - N_G(\text{supp}(x))$ , a contradiction. So  $V = C_1 \cup C_2 \cup C_3$ , done.

Case 2.  $(|C_1|,|C_2|,|C_3|)=(2,1,1)$ :  $v\in N_G(\operatorname{supp}(x))$  and 3-connectivity imply  $|V\setminus\operatorname{supp}(x)\setminus N_G(\operatorname{supp}(x))|\geq 1$ . As  $K_5\not\subseteq\overline{G}-N_G(\operatorname{supp}(x))$ , we know  $|V\setminus\operatorname{supp}(x)\setminus N_G(\operatorname{supp}(x))|=1$ . So  $\overline{G}-N_G(\operatorname{supp}(x))=K_{2,1,1,1}$ . Since  $\overline{G}$  is maximally planar,  $N_G(\operatorname{supp}(x))=\emptyset$ . So  $\overline{G}=K_{2,1,1,1}$  and hence  $\lambda(G)\geq |V(G)|-5$ .

Case 3.  $(|C_1|, |C_2|, |C_3|) = (k, 1, 1), k > 2$ : Now,  $V \setminus \text{supp}(x) \setminus N_G(\text{supp}(x)) = \emptyset$  since otherwise  $K_{3.3} \subseteq \overline{G} - N_G(\text{supp}(x))$ . Also,  $N_G(\text{supp}(x)) = \emptyset$ , since  $v \in N_G(\text{supp}(x))$ 

would be 3-connected to  $C_1$ , which implies  $\overline{G} \ge K_{3,3}$ . So  $V = C_1 \cup C_2 \cup C_3$ , and as  $\overline{G}$  is maximally planar,  $\overline{G}[C_1]$  is a cycle. Hence, any  $v \in C_1$  has degree 4 in  $\overline{G}$ , a contradiction.

Case 4.  $(|C_1|, |C_2|, |C_3|) = (2, 2, 1)$ : then  $V \setminus \text{supp}(x) \setminus N_G(\text{supp}(x)) = \emptyset$ . Hence, the unique vertex  $v \in C_3$  has degree 4 in  $\overline{G}$ , contradiction.

Case 5.  $(|C_1|, |C_2|, |C_3|) = (2, 2, 2)$ : then  $K_{2,2,2} \subseteq \overline{G}$ , contradicting our assumption.

Since  $\ker(M)$  contains no broken vectors, it is a valid representation of G.

The graph G as shown in figure 2 has a planar complement, but  $K_{2,2,2} \subseteq \overline{G}$  and  $\lambda(G) = 1 \not\geq 2 = |V(G)| - 5$ .

## References

- [1] R. Bacher, Y. Colin de Verdière: Multiplicités des valeurs propres et transformations étoile-triangle des graphes, *Bulletin de la Société Mathématique de France*, **123** (1995), 101–117.
- [2] A. E. BROUWER, A. M. COHEN, A. NEUMAIER: Distance-Regular Graphs, Springer Verlag, Ergebnisse der Mathematik und ihrer Grensgebiete 3. Folge, Band 18, 1989.
- [3] Y. Colin de Verdière: On a new graph invariant and a criterion for planarity in: Graph Structure Theory (N. Robertson, P. Seymour, eds.), Contemporary mathematics, American Mathematical Society, Providence, Rhode Island, 1993, 137–147.
- [4] H. VAN DER HOLST: Topological and Spectral Graph Characterisations, Ph.D. Thesis, University of Amsterdam (1996).
- [5] H. VAN DER HOLST, M. LAURENT, and A. SCHRIJVER: On a minor-monotone graph invariant, J. Combin. Theory Ser. B, 65 (1995), 291–304.
- [6] H. VAN DER HOLST, L. LOVÁSZ, A. SCHRIJVER: On the invariance of Colin de Verdière's graph parameter under clique sums, *Linear Algebra and its Applica*tions, 226 (1995), 509–517.
- [7] H. VAN DER HOLST, L. LOVÁSZ, A. SCHRIJVER: The Colin de Verdière graph parameter, to appear in the proceedings of the 16th British Combinatorial Conference.
- [8] A. KOTLOV, L. LOVÁSZ, S. VEMPALA: The Colin de Verdière number and sphere representations of a graph, preprint, 1996.
- [9] L. Lovász and A. Schrijver: The Colin de Verdière number of linklessly embeddable graphs, *Proceedings of the American Mathematical Society*, to appear.

[10] N. ROBERYTSON, P. SEYMOUR, and R. THOMAS: Sachs' Linkless Embedding Conjecture, Combin. Theory Ser. B, 64 (1995), 185–227.

# R. Pendavingh

Faculteit der Wiskunde, Informatica, Natuurkunde en Sterrenkunde Universiteit van Amsterdam Plantage Muidergracht 24 1018 TV Amsterdam